Reusing Information in Multi-Goal Reachability Analyses *

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Abstract: It is known that model checkers can generate test inputs as witnesses for reachability specifications (or, equivalently, as counterexamples for safety properties). While this use of model checkers for testing yields a theoretically sound test-generation procedure, it scales poorly for computing complex test suites for large sets of test goals, because each test goal requires an expensive run of the model checker. We represent test goals as automata and exploit relations between automata in order to reuse existing reachability information for the analysis of subsequent test goals. Exploiting the sharing of sub-automata in a series of reachability queries, we achieve considerable performance improvements over the standard approach. We show the practical use of our multi-goal reachability analysis in a predicate-abstraction-based test-input generator for the test-specification language FQL.

Overview

We consider the problem of performing many reachability queries on a program that is given as source code. Querying a model checker repeatedly for path-sensitive reachability information [BCH\textsuperscript{+}04b] has many interesting applications, e.g., to decompose verification tasks, but most prominently to generate test cases from counterexample paths [BCH\textsuperscript{+}04a, HSTV08]. If, for example, we want to achieve basic-block coverage, we will for each basic block \(b\) try to construct a path through the program that witnesses a program execution that reaches \(b\). In our approach, we describe test-coverage criteria using the coverage-specification language FQL [HSTV08, HSTV09, HSTV10, HTSV10, Hol13], which provides a concise specification of complex coverage criteria. We translate an FQL coverage criterion into a (possibly huge) set of test goals. Each such test goal is represented as a finite automaton, called test-goal automaton, and specifies a reachability query. The model-checking engine then takes a test-goal automaton to restrict the state-space search to the specified paths (for which test cases are desired). Test-goal automata often have identical parts which let us reuse analysis results across several queries. We developed an

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approach that exploits the automaton structure of reachability queries to efficiently reuse reachability results when solving multiple queries. Given two test-goal automata $A$ and $A'$, we introduce the notion of similarity of $A$ and $A'$ modulo a set $X$ of transitions, where $X$ is a subset of the transitions of $A'$. We then identify potentially shared behavior between $A$ and $A'$ via a simulation-modulo-$X$ relation $H$ between the states of $A$ and $A'$. If two states $s$ and $s'$ are related via $H$, then each sequence of $A'$-transitions starting in $s'$ and not including a transition from $X$ corresponds to an equivalent sequence of $A$-transitions starting in $s$. This allows us to reason about feasibility of program executions that are covered by $A'$ based on the reachability results for $A$ as long as we investigate transition sequences shared by both automata [BHTV13].

Because it is generally undecidable whether a test goal is satisfiable on an arbitrary given program, we use an overapproximating reachability analyses, more specifically, a CEGAR-based predicate abstraction [BKW10], to approximate the set of executions of a program until we either (i) have found a partial program execution that is described by a word in the language of the test-goal or (ii) we have shown that there is no such execution. The test-goal automaton guides the reachability analysis, i.e., the analysis tracks program and automaton states simultaneously and stops exploring the state space if there is no possible transition in the program state space or no possible next automaton transition. Based on the excluded transitions $X$, we reuse parts of the already analyzed state space (those parts which do not involve these transitions) or continue state-space exploration along the transitions in $X$. We implemented our approach in the test-input generator CPA/Tiger\(^1\).

References


\(^1\)http://forsyte.at/software/cpatiger/